IN THE CLAIMS:

On amended page 12, at line 1, cancel "New Patent Claims" and substitute --I CLAIM AS MY INVENTION-- therefor;

Amend the following claims 1 through 3.

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1. (Amended) A method [Method] for forming a first commutative checksum [(KP1)] for digital data comprising the steps of [which are grouped into a number of data segments (Di, i = 1 .. n), by a computer,]

grouping said digital data into a number of data segments by a computer.

forming [a) in which] a first segment checksum [(PSi) is formed] for each said data segment [(Di)],

forming said [b) in which the first commutative checksum [(KP1) is formed] by a commutative operation [(\oplus)] on said [the] first segment checksums [(PSi)], and

cryptographically protecting said [c) in which the] first commutative checksum [(KP1) is cryptographically protected] by using a [at least one] cryptographic operation.

2. (Amended) A method [Method] for checking a predetermined cryptographic commutative checksum comprising the steps of: [which is allocated to digital data which are grouped into a number of data segments, by a computer,]

grouping digital data into a number of data segments by a computer, allocating said predetermined cryptographic checksum to said digital

25 data.

subjecting said [a) in which the cryptographic commutative checksum

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[is subjected] to an inverse cryptographic operation to form a first commutative [cryptographic] checksum [(KP1)],

forming [b) in which] a second segment checksum [(PSj) is formed] for each said data segment [(Dj, j = a .. z)],

forming [c) in which] a second commutative checksum [(KP2) is formed] by a commutative operation [(\oplus)] on said [the] second segment checksums [(PSj)], and

checking said [d) in which the] second commutative checksum [(KP2) is checked] for a match with said [the] first commutative checksum [(KP1)].

3. (Amended) A method [Method] for forming and checking a first commutative checksum [(KP1)] for digital data comprising the steps of: [which are grouped into a number of data segments (Di, i = 1 .. n), by a computer,]

grouping said digital data into a number of data segments by a computer,

forming [a) in which] a first segment checksum [(PSi) is formed] for each said data segment [(Di)]

forming said [b)in which the] first commutative checksum [(KP1) is formed] by a commutative operation [(\oplus)] on said first [the] segment checksums [(PSi)],

cryptographically protecting said [c) in which the] first commutative checksum [(KP1) is cryptographically protected] by using at least one cryptographic operation, which forms a cryptographic commutative checksum [being formed],

subjecting said [d) in which the] cryptographic commutative checksum [(KP1) is subjected] to an inverse cryptographic operation to form a reconstructed first [reconstructed] cryptographic checksum [(KP1)],

forming/[e) in which] a second segment checksum [(PSj) is formed] for

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each <u>said</u> data segment [(Dj, j = a ... z)] of <u>said</u> [the] digital data to which <u>said</u> [the] first commutative checksum [(KP1)] is allocated,

forming [f) in which] a second commutative checksum [(KP2) is formed] by a commutative operation [(⊕)] on said [the] second segment checksums [(PSj)], and

checking said [g) in which the] second commutative checksum [(KP2) is checked] for a match with said [the] reconstructed first [reconstructed] commutative checksum (KP1).

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Cancel claim 4 and substitute the following claims 21, 22, and 23

0 therefor.

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21. A method according to claim 1, further comprising the step of:

forming said first segment checksums in accordance with a type selected from the group consisting of a hashing value, a CRC code, and a cryptographic one-way function.

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A method according to claim 2, further comprising the step of: forming said second segment checksums in accordance with a type selected from the group consisting of a hashing value, a CRC code, and a cryptographic one-way function.

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2)
23. A method according to claim 3, further comprising the step of:
forming said first segment checksums and said second segment
checksums in accordance with a type selected from the group consisting of a
hashing value, a CRC code, and a cryptographic one-way function.

Cancel claims 5 and 6, and substitute the following claims 24, 25, and 24 therefor.

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24. A method according to claim/1, wherein:

said cryptographic operation is an operation selected from the group consisting of a symmetric cryptographic method and an asymmetric cryptographic method.

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25. A method according to claim 2, wherein:

said cryptographic operation is an operation selected from the group consisting of a symmetric cryptographic method and an asymmetric cryptographic method.

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26. A method according to claim 3, wherein:

said cryptographic operation is an operation selected from the group consisting of a symmetric cryptographic method and an asymmetric cryptographic method.

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Cancel claim 7 and substitute the following claims 27, 28, and 29 therefor.

25 27. A method according to claim 1, wherein: said commutative operation exhibits the property of associativity.

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26 28. A method according to claim 2, wherein: said commutative operation exhibits the property of associativity.

27. A method according to claim 3, wherein: said commutative operation exhibits the property of associativity.

Cancel claim 8 and substitute the following claims-30, 31, and 32 therefor.

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28 A method according to claim 1, further comprising the step of: 30. protecting said digital data wherein said data segments have no ties to a specific ordering.

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31. A method according to claim 2, further comprising the step of: protecting said digital data wherein said data segments have no ties to a specific ordering.

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32. A method according to claim 3, further comprising the step of: protecting said digital data wherein said data segments have no ties to a specific ordering.

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Cancel claim 9 and substitute the following claims 33, 34, and 35 therefor.

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33. A method according to claim 1, further comprising the steps of: protecting said digital data, and

processing said digital data in accordance with a network management protocol.

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A method according to claim 2, further comprising the steps of: 34. protecting said digital data, and

processing said digital data in accordance with a network management protocol.



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35. A method according to claim 3, further comprising the steps of: protecting said digital data, and

processing said digital data in accordance with a network management

protocol

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Amend the following claims 10 through 12.

10. (Amended) An arrangement [Arrangement] for forming a first commutative checksum [(KP1)] for digital data which are grouped into a number of data segments [(Di, i = 1 .. n)], said arrangement comprising:

[by means of] an arithmetic and logic unit, [which is arranged in such a manner that]

- [a)] a first segment checksum, which [(PSi)] is formed for each said data segment [(Di)],
- [b) the first commutative checksum (KP1) is formed by] a commutative operation [(⊕)] which forms said first commutative checksum by operating on said [the] segment checksums [(Psi)], and
- [c) the first commutative checksum (KP1) is cryptographically protected by using at least one] a cryptographic operation which cryptographically protects said first commutative checksum.
- 11. (Amended) An arrangement [Arrangement] for checking a predetermined first commutative checksum which is allocated to digital data which are grouped into a number of data segments, said arrangement comprising:

[by means of] an arithmetic and logic unit, [which is arranged in such a manner that]

[a) the cryptographic commutative checksum is subjected to] an inverse cryptographic operation to form a first cryptographic checksum [(KP1)] from a

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cryptographic commutative checksum formed by a cryptographic operation,

- [b)] a second segment checksum [(Psj)] which is formed for each said data segment [(Dj, j = a ... z)],
- [c) a second commutative checksum (KP2) is formed by] a commutative operation [(⊕)] which operates on said [the] second segment checksums [(PSj)] which forms a second commutative checksum, and
- [d)] a comparator which checks for a match between said [the] second commutative checksum [(KP2) is checked for a match with the] and said first commutative checksum [(KP1)].
- 12. (Amended) An arrangement [Arrangement] for forming and checking a first commutative checkeum [(KP1)] for digital data which is grouped into a number of data segments [(Di, i = 1 .. n)], said arrangement comprising:

[by means of] an arithmetic and logic unit, [which is arranged in such a manner that]

- [a)] a first segment checksum, which [(PSi)] is formed for each said data segment [(Di)],
- [b) the first commutative checksum (KP1) is formed by] a commutative operation [(\oplus)] which forms said first commutative checksum by operating on said first [the] segment checksums [(Psi)],
- [c) the first commutative checksum (KP1) is cryptographically protected by using at least one] a cryptographic operation which cryptographically protects said first commutative checksum, [a cryptographic commutative checksum being formed,]
- a cryptographic commutative checksum formed by said cryptographic operation.
- [d] the cryptographic commutative checksum is subjected to] an inverse cryptographic operation to form a first cryptographic checksum [(KP1)] from

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said cryptographic commutative checksum,

[e)] a second segment checksum [(PSj)] which is formed for each said data segment [(Dj, j = a z)] of said [the] digital data to which said [the] first commutative checksum [(KP1)] is allocated,

[f) a second commutative checksum (KP2) is formed by] a commutative operation [(#)] which operates on said [the] second segment checksums [(Psj)] which forms a second commutative checksum, and

[g)] a comparator which checks for a match between said [the] second commutative checksum [(KP2) is checked for a match with the] and a reconstructed first [reconstructed] commutative checksum [(KP1)].

Cancel claim 13 and substitute the following claims 36, 37, and 38 therefor.

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36. An arrangement according to claim/10, wherein:

said first segment checksums are formed in accordance with a type selected from the group consisting of a hashing value, a CRC code, and a cryptographic one-way function.

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37. An arrangement according to claim 11, wherein:

said second segment checksums are both formed in accordance with a type selected from the group consisting of a hashing value, a CRC code, and a cryptographic one-way function.

36. An arrangement according to claim 12, wherein:

said first segment checksums and said second segment checksums are both formed in accordance with a type selected from the group consisting of a hashing value, a CRC ode, and a cryptographic one-way function.

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Cancel claims 14 and 15, and substitute the following claims 39, 40, and 41 therefor.

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39. An arrangement according to claim 10 wherein:

said cryptographic operation is an operation/selected from the group consisting of a symmetric cryptographic method and an asymmetric cryptographic method.

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40. An arrangement according to claim 11 wherein:

said cryptographic operation is an operation selected from the group consisting of a symmetric cryptographic method and an asymmetric cryptographic method.

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41. An arrangement according to claim 12 wherein:

said cryptographic operation is an operation selected from the group consisting of a symmetric cryptographic method and an asymmetric cryptographic method.

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Cancel claim 16 and substitute the following claims 42, 43, and 44 therefor.

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An arrangement according to claim 10 wherein said 42. commutative operation exhibits the property of associativity via the arrangement of said arithmetic and logic unit.

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43'. An arrangement according to claim 11 wherein said commutative operation exhibits the property of associativity via the arrangement of said arithmetic and logic unit.

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An arrangement according to claim 12 wherein said commutative operation exhibits the property of associativity via the arrangement of said arithmetic and logic unit.

5 therefor.

Cancel claim 17 and substitute the following claims 45, 46, and 47 r.

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An arrangement according to claim 10 wherein: said digital data are protected, and said data segments have no ties to a specific ordering.

An arrangement according to claim 11 wherein: said digital data are protected, and said data segments have no ties to a specific ordering.

An arrangement according to claim 12 wherein: said digital data are protected, and said data segments have no ties to a specific ordering.

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Cancel claim 18 and substitute the following claims 48, 49, and 50 therefor.

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48. An arrangement according to claim 10 wherein:

said digital data are protected via an arrangement of said arithmetic and logic unit, and

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said digital data are processed in accordance with a network management protocol